Fine-grained Parameterized Algorithms

 $FPA \cdot SoSe\text{-}2024 \cdot \texttt{tcs.uni-frankfurt.de/parameterized/} \cdot 2024\text{-}07\text{-}03 \cdot 2\text{d88037}$



Problem set 7: Longest Path

Overview: With this problem set, you can train reasoning about algebraic methods used in the algorithm for Longest Path.

Instructions: For each skill, select **exactly one** problem below and submit your solution in Moodle; in your submission, make sure to repeat the problem that you are solving. The problems are roughly ordered by difficulty, choose a problem that you find non-trivial and interesting. (You are of course welcome to try the other problems as well and ask us for feedback.)

!! <u>Skill-7a.</u> Longest Path Algorithm: I can fill in the details of the Longest Path algorithm.

7.1 Labels are necessary. In the $2^k n^{O(1)}$ -time algorithm for Longest Path, we introduced labels so that the non-path walks pair up and hence cancel out. This resulted in $O(2^k)$ overhead in the running time. In the $2^{k/2}n^{O(1)}$ -time algorithm for bipartite graphs, we used a different argument for pairing-up non-path walks, by reversing closed subwalks. Does it mean that we can skip the labels then? It seems we cannot, because then we would solve Longest Path in polynomial time! Explain what goes wrong.

7.2 Polynomial Identity Testing. Let p be a polynomial in n variables over the finite field GF(q), such that the degree of p in each variable (i.e., the maximum exponent of a variable) is strictly less than q. Prove that if p evaluates to zero for every argument (i.e., p corresponds to the zero function), then all the coefficients of p are zero (i.e., p is the zero polynomial).

7.3 Multicolored Path. Given a graph *G* and a coloring $c : V(G) \rightarrow [k]$, decide in time $O(2^k \text{poly}(|V|))$ and polynomial space whether *G* contains a *k*-path with vertices from all *k* color classes.

!! <u>Skill-7b.</u> Adapt the Longest Path Algorithm: I can adapt the ideas of the Longest Path algorithm to solve other problems.

7.4 Weighted Longest Path. In the Weighted Longest Path problem, we are given a directed graph G with a weight function $w: E(G) \to \{0, ..., W\}$ and the goal is to find a k-path P of the smallest weight $\sum_{uv \in E(P)} w(u, v)$. Describe a bounded-error randomized algorithm for this problem that runs in time $2^k \cdot W \cdot n^{O(1)}$ and uses polynomial space.

7.5 Triangle Packing. Describe a bounded-error randomized algorithm for the Triangle Packing problem: given a graph G and a number $k \in \mathbb{N}$, decide whether G contains k disjoint triangles (subgraphs isomorphic to K_3). Your algorithm should use polynomial space and run in time $2^{3k}n^{O(1)}$.